

MidTerm

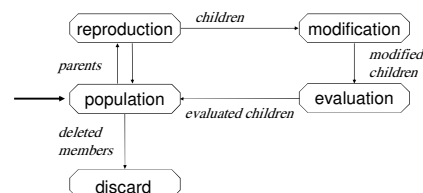
- Q1) Under what circumstances would it make sense to go “down hill” (i.e. explore a child whose heuristic value is worse than the parent’s) when executing the hill-climbing search algorithm? When would it *not* make sense?

Question 2

- Explain how does the Genetic Algorithm work? Why is it considered as Local Search Algorithm
- And Explain!!!
- local search algorithms are useful for solving pure **optimization problems**, in which the aim is to find the **objective function**.

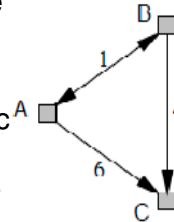
In such cases, can use iterative improvement algorithms; keep a single “current” state, try to improve it

The GA Cycle of Reproduction



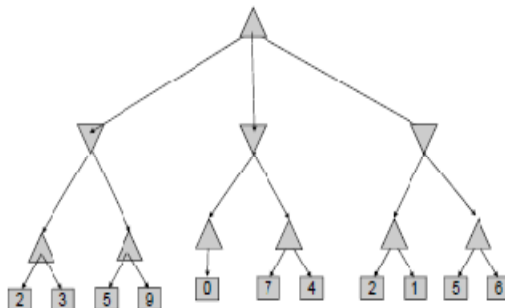
Question 3

- The following diagram represents the state space of a deterministic problem, with each arrow denoting a possible operator (labeled with the step cost). Assume that the successors of a state are generated in alphabetical order, and that there is no repeated state checking.
- Show the search tree generated by breadth-first search applied to the problem of starting in A, where C is the goal. Circle the tree node that the search identifies as the solution.
- What is the branching factor here?
- Which of the following algorithms will find solution in this case breadth first, depth first?



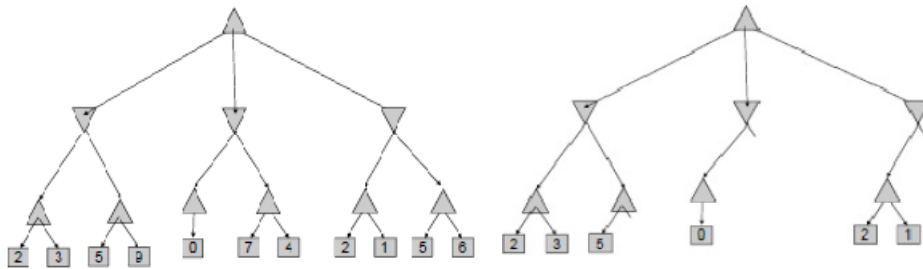
Question 4

- Apply the minimax algorithm to the game tree below, where it is the maximizer's turn to play. The values of the evaluation function of the leaves are listed.
- Write the values of the intermediate nodes

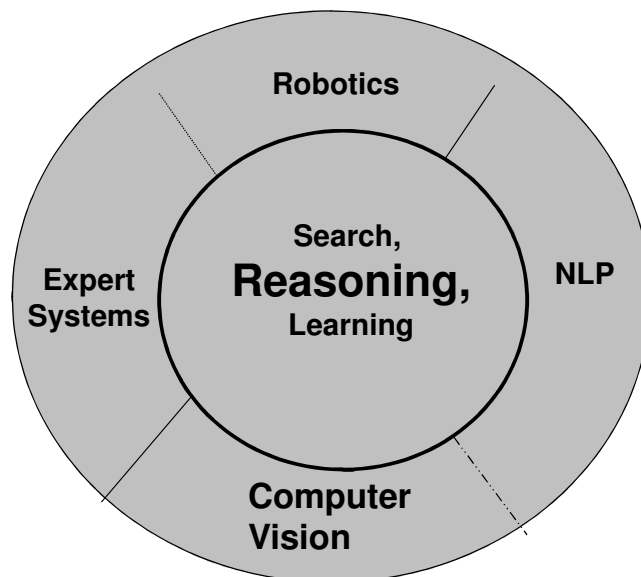


Question 4

- Indicate (mark the edge) the proper move of the maximizer by marking one of the root's outgoing arcs.

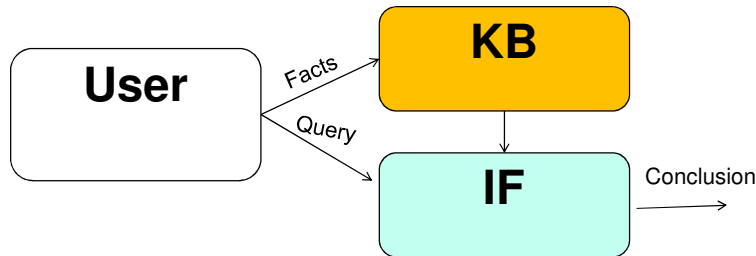


Reasoning



First-order logic Chapter 8-Russel Representation and Reasoning

- In order to determine appropriate actions to take, an intelligent system needs to represent information about the world and draw conclusions based on general world knowledge and specific facts.
- Knowledge is represented by sentences in some language stored in a **knowledge base (KB)**.
- A system draws conclusions from the KB to answer questions, take actions using **Inference Engine (IF)**.



Knowledge Representation

- Logics are formal languages for representing information such that conclusions can be drawn
- **Syntax:** defines the sentences in the language
- **Semantics:** define the “meaning” of sentences: i.e., define truth of a sentence in a world
- E.g., the language of arithmetic
 - $x+2 \geq y$ is a sentence; $x+2 > \{ \}$ is not a sentence
syntax
 - $x+2 \geq y$ is true in a world where $x = 7, y = 1$
 - $x+2 \geq y$ is false in a world where $x = 0, y = 6$ } semantics

Inference

- Logical Inference (deduction) derives new sentences in the language from existing ones,.
Socrates is a man.
All men are mortal.
Socrates is mortal.
- Proper inference should only derive sound conclusions

Logics

- Logics are formal languages for representing information such that conclusions can be drawn
- **Syntax:** defines the sentences in the language
- **Semantics:** define the “meaning” of sentences: i.e., define true of a sentence in a world

Examples of Types of Logics

Language	What exist	Degree of belief of an Agent
Propositional Logic	Facts	{0,1} T or F
First Order Logic	Facts, Objects, Relations	{0,1} T or F
Temporal Logic	Facts, Objects, Relations, Time	{0,1} T or F
Probability Theory	Facts	Chances of belief [0,1]
Fuzzy Logic	Degree of truth about Facts	Degree of belief [0,1]

Propositional calculus & First-order logic

- **Propositional logic** assumes world contains **facts**.
- **First-order logic** (like natural language) assumes the world contains
 - Objects: people, houses, numbers, ...
 - Relations: red, round, prime,...
 - Functions: fatherof, friend, in,...
- **Propositional calculus**
 $A \wedge B \Rightarrow C$
- **First-order predicate calculus**
 $(\forall x)(\exists y) \text{ Mother}(y,x)$

Syntax of PC_{Chapter 7-Russel}

- Connectives: $\neg, \wedge, \vee, \Rightarrow$
- Propositional symbols, e.g., P, Q, R, ...
 - *True, False*
 - **Syntax of PC**
- sentence \rightarrow atomic sentence | complex sentence
- atomic sentence \rightarrow Propositional symbol, *True, False*
- Complex sentence \rightarrow
 - \neg sentence
 - $(\text{sentence} \wedge \text{sentence})$
 - $(\text{sentence} \vee \text{sentence})$
 - $(\text{sentence} \Rightarrow \text{sentence})$
- **Rules of Inference:**
- Ex: Modus ponens

Sentence in PC

A **sentence (also called a formula or well-formed formula or wff)** is defined as:

- A symbol (S, P, ...etc)
- If S is a sentence, then $\neg S$ is a sentence, where " \neg " is the "not" logical operator
- If S and T are sentences, then $(S \vee T)$, $(S \wedge T)$, $(S \Rightarrow T)$, and $(S \Leftrightarrow T)$ are sentences, where the four logical connectives correspond to "or," "and," "implies," and "if and only if," respectively

Example

P means "It is hot"

Q means "It is humid"

R means "It is raining"

Examples of PL sentences:

$(P \wedge Q) \Rightarrow R$ (here meaning "If it is hot and humid, then it is raining")

$Q \Rightarrow P$ (here meaning "If it is humid, then it is hot")

$\neg Q$ (here meaning "It is not humid.")

Semantics of PC

A	B	$\neg A$	$A \wedge B$	$A \vee B$	$A \Rightarrow B$
<i>True</i>	<i>True</i>	<i>False</i>	<i>True</i>	<i>True</i>	<i>True</i>
<i>True</i>	<i>False</i>	<i>False</i>	<i>False</i>	<i>True</i>	<i>False</i>
<i>False</i>	<i>False</i>	<i>True</i>	<i>False</i>	<i>False</i>	<i>True</i>
<i>False</i>	<i>True</i>	<i>True</i>	<i>False</i>	<i>True</i>	<i>True</i>

Semantics of PC

Validity and Inference

- **interpretation of the sentence:** Given the truth values of all of the constituent symbols in a sentence, that sentence can be "evaluated" to determine its truth value (True or False).
- A **model** is an interpretation (i.e., an assignment of truth values to symbols) of a set of sentences such that each sentence is True. A model is just a formal mathematical structure for the world.
- A **valid sentence (also called a tautology)** is a sentence that is True under *all interpretations*. *Hence, no matter what the world is actually like or what the semantics is, the sentence is True.*
For example "It's raining or it's not raining."
Remark: Validity can be checked by the truth table

Semantics of PC

Validity and Inference

- An **inconsistent sentence (also called unsatisfiable or a contradiction)** is a sentence that is False under *all interpretations*.
For example, "It's raining and it's not raining."
- Sentence P **entails sentence Q**, written $P \models Q$, means that whenever P is True, so is Q. In other words, all models of P are also models of Q

Satisfiability

- A sentence is **satisfiable** if it is true under some interpretation (i.e. it has a model), otherwise the sentence is **unsatisfiable**.
- A sentence is **valid** if and only if its negation is unsatisfiable.
- Therefore, algorithms for either validity or satisfiability checking are useful for logical inference.
- If there are n propositional symbols in a sentence, then we must check 2^n rows for validity
- **Satisfiability is NP-complete**, i.e. there is no polynomial-time algorithm to solve.
- Yet, many problems can be solved very quickly.

Rules of Inference

- A sequence of inference rule applications that leads to a desired conclusion is called a **logical proof**.
- $A \vdash B$, denotes that B can be derived by some inference procedure from the set of sentences A.
- Inference rules can be verified by the truth-table
- The truth table method of inference is **complete for PL**
- Then used to construct **sound proofs**.
- Finding a proof is simply a **search problem** with the inference rules as operators and the conclusion as the goal

Rules of Inference

- **Modus Ponens:** $\{\alpha \Rightarrow \beta, \alpha\} \vdash \beta$
- **And Elimination:** $\{\alpha \wedge \beta\} \vdash \alpha$; $\{\alpha \wedge \beta\} \vdash \beta$
- **Double negation Elimination:** $\{\neg\neg\alpha\} \vdash \alpha$
- **Implication Elimination** $\{\alpha \Rightarrow \beta\} \vdash \neg\alpha \vee \beta$
- **Unit resolution:** $\{\alpha \vee \beta, \neg\beta\} \vdash \alpha$
- **Resolution:** $\{\alpha \vee \beta, \neg\beta \vee \gamma\} \vdash \alpha \vee \gamma$

Famous logical equivalences

- $(a \vee b) \equiv (b \vee a)$ *commutativity*
- $(a \wedge b) \equiv (b \wedge a)$ *commutativity*
- $((a \wedge b) \wedge c) \equiv (a \wedge (b \wedge c))$ *associativity*
- $((a \vee b) \vee c) \equiv (a \vee (b \vee c))$ *associativity*
- $\neg(\neg(a)) \equiv a$ *double-negation elimination*
- $(a \Rightarrow b) \equiv (\neg(b) \Rightarrow \neg(a))$ *contraposition*
- $(a \Rightarrow b) \equiv (\neg(a) \vee b)$ *implication elimination*
- $\neg(a \wedge b) \equiv (\neg(a) \vee \neg(b))$ *De Morgan*
- $\neg(a \vee b) \equiv (\neg(a) \wedge \neg(b))$ *De Morgan*
- $(a \wedge (b \vee c)) \equiv ((a \wedge b) \vee (a \wedge c))$ *distributivity*
- $(a \vee (b \wedge c)) \equiv ((a \vee b) \wedge (a \vee c))$ *distributivity*

Pros and cons of propositional logic

- ✓ Propositional logic is declarative:
pieces of syntax correspond to facts
- ✓ Propositional logic is compositional:
meaning of $A \wedge B$ is derived from meaning of A and B
- ✓ Meaning in propositional logic is context-independent
 - (unlike natural language, where meaning depends on context)
- Propositional logic has very limited expressive power
 - (unlike natural language)

Propositional logic is a weak language

- Hard to identify “individuals.” Ex. Mary, 3
- Can’t directly talk about properties of individuals or relations between individuals. Ex. “Bill is tall”
- Generalizations, patterns, regularities can’t easily be represented. Ex. all triangles have 3 sides
- **First-Order Logic** (abbreviated FOL or FOPC) is expressive enough to concisely represent this kind of situation.
 - FOL adds relations, variables, and quantifiers, e.g.,
 - “Every elephant is gray”: $\forall x (\text{elephant}(x) \rightarrow \text{gray}(x))$
 - “There is a white elephant”: $\exists x (\text{elephant}(x) \wedge \text{white}(x))$